

Lean Bourgain Extractor

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Chapter 1

Additive Combinatorics

Lemma 1.1. For any two sets A, B , we have $|A - B| \leq \frac{|A+B|^3}{|A||B|}$

Proof. By the triangle inequality we have $|A - B| \leq \frac{|A+B||B+B|}{|B|}$, and from the Plünnecke-Ruzsa inequality we have $|B + B| \leq \left(\frac{|A+B|}{|A|}\right)^2 |A|$. \square

Lemma 1.2. For any set A and a non-zero value x , we have $|xA| = |A|$

Proof. This is obvious from the bijection of multiplication by x . \square

Lemma 1.3. we have $-(A \cap B) = -A \cap -B$

Proof. Trivial from the definitions. \square

Lemma 1.4. For any set A and two values a, b , we have $(a + b)A \subseteq aA + bA$.

Proof. For any value $(a + b)v$ with $v \in A$ we have $av \in aA$, $bv \in bA$, and $av + bv = (a + b)v$ \square

Lemma 1.5. For any set A and two values a, b , we have $(a - b)A \subseteq aA - bA$.

Proof. Exactly the same as the previous lemma. \square

Lemma 1.6. If $A \cap C \neq \emptyset$, we have $|B + C| \leq \frac{|B+A||C+C|}{|A \cap C|}$.

Proof. By the triangle inequality, we have $|B + C| \leq \frac{|B+(A \cap C)|| (A \cap C) + C|}{|A \cap C|}$, and this is less than $\frac{|B+A||C+C|}{|A \cap C|}$ because $B + (A \cap C) \subseteq B + A$ and $(A \cap C) + C \subseteq C + C$. \square

Lemma 1.7. For any three sets A, B, C , we have $|A + B + C| \leq \frac{|C+A||A+B|^8}{|A|^6|B|^2}$

Proof. If either A or B are empty this is trivial. Otherwise we have an element $v \in B$. We obviously have $(A + B) \cap (A + \{v\}) = A + \{v\}$, and it is nonempty. So by 1.6 we have

$$|A+B+C| = |C+(A+B)| \leq \frac{|C+A+\{v\}|| (A+B) + (A+B) |}{|A+\{v\}|} = \frac{|C+A|| (A+B) + (A+B) |}{|A|}$$

By Ruzsa's covering lemma we have a set u of size $\leq \frac{|A+B|}{|B|}$ such that $A \subseteq u + B - B$. This gives $\frac{|C+A|| (A+B) + (A+B) |}{|A|} \leq \frac{|C+A|| (u+B-B+B) + (u+B-B+B) |}{|A|} = \frac{|C+A|| 2 \cdot u + 4 \cdot B - 2 \cdot B |}{|A|} \leq \frac{|C+A|| u^2 | 4 \cdot B - 2 \cdot B |}{|A|}$.

By the Plünnecke-Ruzsa inequality we now have $|4 \cdot B - 2 \cdot B| \leq \left(\frac{|A+B|}{|A|}\right)^6 |A|$, and the result follows from this and the bound on $|u|$. \square

Lemma 1.8. *We have that $Q(A, xA)$ for $x \neq 0$ is the number of quadruples $(a, b, c, d) \in A^4$ such that $a + xb = c + xd$.*

Proof. By a direct bijection from quadruples $(a, b, c, d) \in A \times (xA) \times A \times (xA)$ such that $a + b = c + d$. \square

Chapter 2

Growth

Theorem 2.1. *For any set A over a finite field of size q there is a value $a \neq 0$ such that $|A + aA| \geq \frac{\min(|A|^2, q)}{2}$.*

Proof. First, we show it's sufficient for a to have $Q(A, aA) \leq |A|^2 + \frac{|A|^2(|A|^2-1)}{q-1}$. We have $|A + aA| \geq \frac{|A|^2|aA|^2}{Q(A, aA)} \geq \frac{|A|^4}{|A|^2 + \frac{|A|^2(|A|^2-1)}{q-1}}$. We need to show $\frac{x^2}{x + \frac{x(x-1)}{q-1}} \geq \frac{\min(x, q)}{2}$. If $x < q$ then

$$\frac{x^2}{x + \frac{x(x-1)}{q-1}} \geq \frac{x^2}{x + \frac{x}{2}} = \frac{x}{2}.$$

Otherwise, if $q \leq x$, we need to show

$$\frac{x^2}{x + \frac{x(x-1)}{q-1}} - \frac{q}{2} \geq 0.$$

By directly expanding, we have

$$\frac{x^2}{x + \frac{x(x-1)}{q-1}} - \frac{q}{2} = \frac{(q-2)(x-q)}{2(q+x-2)}$$

We have $2 \leq q$, so this value is nonnegative. Now to show that for some a , $Q(A, aA) \leq |A|^2 + \frac{|A|^2(|A|^2-1)}{q-1}$. Now we show it suffices to show that $\sum_{a \neq 0} Q(A, aA) \leq |A|^2(q-1) + |A|^2(|A|^2-1)$.

This is because if all values were larger than $|A|^2 + \frac{|A|^2(|A|^2-1)}{q-1}$, the sum couldn't've been so small.

To show $\sum_{a \neq 0} Q(A, aA) \leq |A|^2(q-1) + |A|^2(|A|^2-1)$, we can use 1.8. The quadruples with $a = c, b = d$ contribute at most $|A|^2(q-1)$, and the other quadruples contribute at most $|A|^2(|A|^2-1)$, because they determine a unique a . \square

Theorem 2.2. *For any set A in \mathbb{F}_p (for p prime), we have $|3A^2 - 3A| \geq \frac{\min(|A|^2, p)}{2}$.*

Proof. If $|A| \leq 1$ then $|3A^2 - 3A| = |A|$, and this is greater than $|A|^2/2$. Otherwise, we split into cases by whether $\frac{A-A}{A-A}$ is the entire universe. If it is, then by 2.1 we have some value $v = (a-b)/(c-d)$ such that $|A+vA| \geq \frac{\min(|A|^2, p)}{2}$. By 1.2, we have $|A+vA| = |(c-d)A + (a-b)A|$.

Now $(c-d)A + (a-b)A$, by 1.5, this is a subset of $cA + aA - dA - bA$, which is a subset of $2A^2 - 2A^2$, and then $|3A^2 - 3A^2| = |A^2 - A^2 + 2A^2 - 2A^2| \geq |2A^2 - 2A^2| \geq \frac{\min(|A|^2, p)}{2}$.

Otherwise, there must be some value such that $v = (a-b)/(c-d)$ such that $(a-b+c-d)/(c-d) = (a-b)/(c-d) + 1 \notin \frac{A-A}{A-A}$. Because of that $|A + (a-b+c-d)/(c-d)A| = |A|^2$, so $|(c-d)A + (a-b+c-d)A| = |A|^2$. Using 1.5 and 1.4 we have $(c-d)A + (a-b+c-d)A \subseteq 3A^2 - 3A^2$, so $|3A^2 - 3A^2| \geq |A|^2 \geq \frac{|A|^2}{2}$. \square

Chapter 3

Stabilizer

Definition 3.1. $\text{Stab}_K(A)$ is the set $\{x \mid |A + xA| \leq K|A|\}$.

Lemma 3.2. For $a \in \text{Stab}_K(A)$, we also have $a^{-1} \in \text{Stab}_K(A)$.

Proof. If $a = 0$ this is trivial, otherwise by 1.2 we have $|A + a^{-1}A| = |a(A + a^{-1}A)| = |A + aA| \leq K|A|$. \square

Lemma 3.3. If $A \neq \emptyset$ and $a \in \text{Stab}_K(A)$ then $1 \leq K$.

Proof. Trivial. \square

Lemma 3.4. For $a \in \text{Stab}_K(A)$, we also have $a \in \text{Stab}_{K^3}(A)$.

Proof. If $a = 0$ or $A = \emptyset$ this is trivial, otherwise by 1.1 and 1.2 we have

$$|A - aA| \leq \frac{|A + aA|^3}{|A||aA|} = \frac{|A + aA|^3}{|A|^2} \leq \frac{K^3|A|^3}{|A|^2} = K^3|A|$$

\square

Lemma 3.5. We have $-\text{Stab}_K(A) \subseteq \text{Stab}_{K^3}(A)$.

Proof. Immediate from 3.4. \square

Lemma 3.6. For $a \in \text{Stab}_{K_1}(A), b \in \text{Stab}_{K_2}(A)$, we have $a + b \in \text{Stab}_{K_1^8 K_2}(A)$.

Proof. If $a = 0$ or $A = \emptyset$ this is trivial. Otherwise, we have $A + (a + b)A \subseteq A + aA + bA$, by 1.4, and by 1.7 and 1.2 we have $|A + aA + bA| \leq \frac{|A + bA||A + aA|^8}{|A|^8} \leq K_1^8 K_2 |A|$. \square

Lemma 3.7. We have $\text{Stab}_{K_1}(A) + \text{Stab}_{K_2}(A) \subseteq \text{Stab}_{K_1^8 K_2}(A)$.

Proof. Immediate from 3.6. \square

Lemma 3.8. For $n \in \mathbb{N}$ we have $(n + 1) \cdot \text{Stab}_K(A) \subseteq \text{Stab}_{K^{8n+1}}(A)$.

Proof. By induction with 3.7. \square

Lemma 3.9. We have $\text{Stab}_{K_1}(A) - \text{Stab}_{K_2}(A) \subseteq \text{Stab}_{K_1^8 K_2^3}(A)$.

Proof. Immediate from 3.7 and 3.5. \square

Lemma 3.10. For $a \in \text{Stab}_{K_1}(A), b \in \text{Stab}_{K_2}(A)$, we have $ab \in \text{Stab}_{K_1K_2}(A)$.

Proof. If $a = 0$ this is trivial with 3.3. Otherwise, we have, by 1.2 $|A + abA| = |a^{-1}A + bA|$. By the triangle inequality we have $|a^{-1}A + bA| \leq \frac{|A+a^{-1}A||A+bA|}{|A|}$, and using 3.2 we get that this is $\leq K_1K_2|A|$. \square

Lemma 3.11. We have $\text{Stab}_{K_1}(A)\text{Stab}_{K_2}(A) \subseteq \text{Stab}_{K_1K_2}(A)$.

Proof. Immediate from 3.10. \square

Lemma 3.12. If $a \in \text{Stab}_K(A)$ and $K \leq K'$ then $a \in \text{Stab}_{K'}(A)$.

Proof. Trivial from the definition. \square

Lemma 3.13. If $K \leq K'$ then $\text{Stab}_K(A) \subseteq \text{Stab}_{K'}(A)$.

Proof. Trivial from 3.12 \square

Lemma 3.14. If $1 \leq K$ implies $K \leq K'$ then $\text{Stab}_K(A) \subseteq \text{Stab}_{K'}(A)$.

Proof. If $A = \emptyset$ this is trivial. Otherwise, if $K < 1$ then from 3.3 $\text{Stab}_K(A) = \emptyset$ and this is trivial. Otherwise we get 3.13. \square

Lemma 3.15. We have $3\text{Stab}_K(A)^2 - 3\text{Stab}_K(A)^2 \subseteq \text{Stab}_{K^{374}}(A)$.

Proof. Immediate from 3.8, 3.9 and 3.11. \square

Lemma 3.16. We have $\frac{\min(|\text{Stab}_K(A)|^2, p)}{2} \leq |\text{Stab}_{K^{374}}(A)|$.

Proof. Immediate from 3.15 and 2.2. \square

Lemma 3.17. If $4 \leq |\text{Stab}_K(A)|$, then $\min(|\text{Stab}_K(A)|^{\frac{3}{2}}, \frac{p}{2}) \leq |\text{Stab}_{K^{374}}(A)|$.

Proof. From direct calculation using 3.16. \square

Lemma 3.18. If $4 \leq |\text{Stab}_K(A)|$ for all $n \in \mathbb{N}$, $\min(|\text{Stab}_K(A)|^{(\frac{3}{2})^n}, \frac{p}{2}) \leq |\text{Stab}_{K^{374^n}}(A)|$.

Proof. By induction on 3.17. \square

Definition 3.19. $\text{StabC}_2(\beta) = 374^{\lceil \log_{\frac{3}{2}}(1/\beta) \rceil}$

Lemma 3.20. If $4 \leq |\text{Stab}_K(A)|$ and $p^\beta \leq |\text{Stab}_K(A)|$, then $\frac{p}{2} \leq |\text{Stab}_{K^{\text{StabC}_2(\beta)}}(A)|$.

Proof. By setting $n = \text{StabC}_2(\beta)$ at 3.18. \square

Definition 3.21. $\text{StabC}(\beta) = 9\text{StabC}_2(\beta)$

Lemma 3.22. If $4 \leq |\text{Stab}_K(A)|$ and $p^\beta \leq |\text{Stab}_K(A)|$, then $\text{Stab}_{K^{\text{StabC}(\beta)}}(A) = \mathbb{F}$.

Proof. By Cauchy-Davenport and 3.7 after ???. \square

Lemma 3.23. If $p^\beta \leq |A| \leq p^{1-\beta}$ and $K < \frac{p^\beta}{2}$, then $\text{Stab}_K(A) \neq \mathbb{F}$.

Proof. 2.1 gives a value a which by direct computation we can show isn't in $\text{Stab}_K(A)$. \square

Lemma 3.24. If $4 \leq |\text{Stab}_K(A)|, p^\beta \leq |\text{Stab}_K(A)|, p^\gamma \leq |A| \leq p^{1-\gamma}$ then $\frac{p^\gamma}{2} \leq K^{\text{StabC}(\beta)}$

Proof. By applying 3.22 and 3.23. \square

Chapter 4

Energy Growth

Theorem 4.1. *Let $S_1, S_2, \dots, S_k \subseteq S$ be finite sets with $|S_i| \geq \delta|S|$ for all i . Then, there exists i such that $|\{j \mid |S_j \cap S_i| \geq (\delta^2/2)k\}| \geq (\delta^2/2)k$*

Proof. This is exactly Claim 3.3.6 in [Dvi12]. □

Theorem 4.2. *Let A, T be finite sets with $Q(A, \lambda A) \geq \frac{|A|^3}{K}$ for all $\lambda \in T$. Then there exist sets A', T' with $\frac{|A|}{16K} \leq |A'|$ and $\frac{|T|}{2^{17}K^4} \leq |T'|$, such that $T' \subseteq \text{Stab}_{2^{110}K^{42}}(A')$.*

Proof. This is exactly Theorem 3.3.5 in [Dvi12], using BSG from LeanAPAP. □

[Dvi12]: Dvir, Zeev. Incidence Theorems and Their Applications, now, 2012, doi: 10.1561/04000000056.

Chapter 5

Lines

TOOD: Figure out how to write blueprints about definitions

Definition 5.1. A line over a field \mathbb{F} is a linear subspace of \mathbb{F}^3 of dimension 2.

Definition 5.2. A point $(x, y) \in \mathbb{F}^2$ is in a line L iff $(x, y, 1) \in L$.

Definition 5.3. Given a linear isomorphism P and a line L , we have a line PL .

Proof. This is a valid line because linear isomorphism preserves dimension. □

Lemma 5.4. For any linear equivalence, applying it to lines is injective.

Proof. From the injectivity of linear isomorphisms. □

Theorem 5.5. Given a set P of points and a set L of lines, the number of incidences is at most $\sqrt{|L||P|(|P| + |L|)}$.

TODO2

Chapter 6

Projective Transformations

TOOD: Figure out how to write blueprints about definitions

Definition 6.1. *Given two different values, $(x_1, y_1), (x_2, y_2) \in \mathbb{F}^2$, $(x_1, y_1) \neq (x_2, y_2)$, we get a linear isomorphism A such that $A(x_1, y_1, 1) = (1, 0, 0)$ and $A(x_2, y_2, 1) = (0, 1, 0)$.*

Lemma 6.2. *Given a point p not on the line between $(x_1, y_1), (x_2, y_2)$, the projective transformation defined by those points doesn't move it to infinity.*

Proof. Direct calculation. □

Chapter 7

Incidence

Definition 7.1. $C = C_2 + 1$

Definition 7.2. $\varepsilon(\beta) = \varepsilon_2(\beta)/3$

Theorem 7.3. *Let there be a set P of points and a set L of lines over a prime field, with $|P| \leq n, |L| \leq n$ and $p^\beta \leq n \leq p^{2-\beta}$. Then the number of intersections is at most $Cn^{\frac{3}{2}-\varepsilon(\beta)}$.*

Proof. We reduce this to 7.5, by removing all points contained in at most $n^{\frac{1}{2}-\varepsilon(\beta)}$ lines. This removes at most $n^{\frac{3}{2}-\varepsilon(\beta)}$ points, which is corrected for with $C = C_2 + 1$. \square

Definition 7.4. $C_2 = \sqrt{2(C_3 + \frac{\sqrt{2}}{4})}$

Theorem 7.5. *Let there be a set P of points and a set L of lines over a prime field, with $|P| \leq n, |L| \leq n$ and $p^\beta \leq n \leq p^{2-\beta}$, and each point intersecting with at least $n^{\frac{1}{2}-\varepsilon(\beta)}$ lines. Then the number of intersections is at most $C_2n^{\frac{3}{2}-\varepsilon(\beta)}$.*

Proof. We reduce this to 7.6, by removing all points contained in more than $4n^{\frac{1}{2}+\varepsilon(\beta)}$ lines. There can be at most $n^{1-2\varepsilon(\beta)} \frac{\sqrt{2}}{4}$ such points, by 5.5. Therefore, there are still many remaining points, and because each point has at least $n^{\frac{1}{2}-\varepsilon(\beta)}$ lines there are still many intersections. \square

Theorem 7.6. *Let there be a set P of points and a set L of lines over a prime field, with $|P| \leq n, |L| \leq n$ and $p^\beta \leq n \leq p^{2-\beta}$, and each point contained in at least $n^{\frac{1}{2}-\varepsilon(\beta)}$ lines and at most $4n^{\frac{1}{2}+\varepsilon(\beta)}$. Then the number of intersections is at most $C_3n^{\frac{3}{2}-\varepsilon_2(\beta)}$.*

Proof. We use ?? to claim that there exist two points, a, b such that for a large number of points they are both on a line from a and a line from b . We only keep those, and because all points are contained in $n^{\frac{1}{2}-\varepsilon(\beta)}$ lines there are still many intersections. Then we remove all points on the line between a and b . Because all lines, except maybe one, intersect at most one such point, this step doesn't remove many intersections. Now we can apply 7.7. \square

Theorem 7.7. *Let there be a set P of points and a set L of lines over a prime field, with $|P| \leq n, |L| \leq n$ and $p^\beta \leq n \leq p^{2-\beta}$, two different points p_1, p_2 , which are both contained in at most $4n^{\frac{1}{2}+\varepsilon(\beta)}$ lines, with no points in P on the line p_1p_2 , and all points in P on an intersection of some line from p_1 and some line from p_2 . Then the number of intersections is at most $C'n^{\frac{3}{2}-\varepsilon'(\beta)}$.*

Proof. By 6.1 we can reduce this to 7.8. TODO. \square

Theorem 7.8. *Let there be two sets A, B and a set L of lines over a prime field, with $|A| \leq 4n^{\frac{1}{2}+2\epsilon(\beta)}$, $|B| \leq 4n^{\frac{1}{2}+2\epsilon(\beta)}$, $|L| \leq n$ and $p^\beta \leq n \leq p^{2-\beta}$. Then the number of intersections is at most $C'n^{\frac{3}{2}-\epsilon'(\beta)}$.*

Proof. We reduce to 7.9 by removing all lines which contain too few points. \square

Theorem 7.9. *Let there be two sets A, B and a set L of lines over a prime field, with $|A| \leq 4n^{\frac{1}{2}+2\epsilon(\beta)}$, $|B| \leq 4n^{\frac{1}{2}+2\epsilon(\beta)}$, $|L| \leq n$ and $p^\beta \leq n \leq p^{2-\beta}$. Additionally, suppose there are at least $n^{\frac{1}{2}-\epsilon'(\beta)}$ points on each line. Then the number of intersections is at most $C'_2 n^{\frac{3}{2}-\epsilon'(\beta)}$.*

Proof. We now remove all horizontal lines, to reduce to 7.10. This doesn't remove many intersections because each point can intersect at most one horizontal line. \square

Theorem 7.10. *Let there be two sets A, B and a set L of non-horizontal lines over a prime field, with $|A| \leq 4n^{\frac{1}{2}+2\epsilon(\beta)}$, $|B| \leq 4n^{\frac{1}{2}+2\epsilon(\beta)}$, $|L| \leq n$ and $p^\beta \leq n \leq p^{2-\beta}$. Additionally, suppose there are at least $n^{\frac{1}{2}-\epsilon'(\beta)}$ points on each line. Then the number of intersections is at most $C'_2 n^{\frac{3}{2}-\epsilon'(\beta)}$.*

Proof. We apply ?? to get two values $b_1, b_2 \in B$ such that many lines pass through these rows. Because there are many points on each line, only keeping those still gives many incidences. Now, a line can be described as two points a_1, a_2 , and the line would be the line passing through $(a_1, b_1), (a_2, b_2)$. Suppose it passes through a given point (a, b) . This gives $a = \frac{b_2-b}{b_2-b_1}a_1 + \frac{b-b_1}{b_2-b_1}a_2$, so $\frac{b_2-b}{b_2-b_1}a_1 + \frac{b-b_1}{b_2-b_1}a_2 \in A$. Equivalently, there are many $(N^{3/2-\epsilon})$ triplets $(b, a_1, a_2) \in B \times A \times A$ such that $\frac{b_2-b}{b_2-b_1}a_1 + \frac{b-b_1}{b_2-b_1}a_2 \in A$. This implies that there must be many $(N^{1/2-\epsilon})$ values of b such that there is a large number of pairs (a_1, a_2) with this property. Now we can only keep those, remove b_1, b_2 , and apply 7.11 \square

Theorem 7.11. *Let there be two sets A, B and a set L of non-horizontal lines over a prime field, with $|A| \leq 4n^{\frac{1}{2}+2\epsilon(\beta)}$, $|B| \leq 4n^{\frac{1}{2}+2\epsilon(\beta)}$, $|L| \leq n$ and $p^\beta \leq n \leq p^{2-\beta}$. Suppose there are at least $n^{\frac{1}{2}-\epsilon'(\beta)}$ points on each line, and lastly, suppose there are two values $b_1, b_2 \notin B$, TODO. Then $|B|$ is at most $C'_5 n^{1/2-\epsilon'(\beta)-\epsilon'(\beta)-4\epsilon(\beta)}$.*

Proof. TODO \square

Chapter 8

Transfer operator

Definition 8.1. For $f : A \rightarrow B, G : A \rightarrow C$ we have $f\#g$ is a function $B \rightarrow C$ defined by $f\#g(x) = \sum_{f(y)=x} g(y)$.

Proposition 8.2. We have $f\#(g + h) = f\#g + f\#h$.

Proposition 8.3. We have $f\#(g - h) = f\#g - f\#h$.

Proposition 8.4. If h is an additive homomorphism we have $h \circ (f\#g) = f\#(g \circ h)$.

Proposition 8.5. If f is a bijection we have $(f\#g)(x) = g(f^{-1}(x))$.

Lemma 8.6. We have $h\#(f\#g) = (h \circ f)\#g$.

Proof.

$$\sum_{y \in h^{-1}(x)} \sum_{z \in f^{-1}(y)} g(z) = \sum_z \sum_{y \in h^{-1}(x), z \in f^{-1}(y)} g(z) = \sum_z [h(f(z)) = x] g(z) = \sum_{z \in (h \circ f)^{-1}(x)} g(z) = ((h \circ f)\#g)(x)$$

□

Proposition 8.7. $\text{id}\#f = f$

Proposition 8.8.

$$\sum_x (f\#g)(x)h(x) = \sum_x g(x)h(f(x))$$

Lemma 8.9.

$$E[(f\#g)(x)h(x)] = \frac{|A|}{|B|} E[g(x)h(f(x))]$$

Proof. By unfolding the expectation and using 8.8. □

Proposition 8.10. if $(f\#g)(x) \neq 0$ then $\exists y, f(y) = x$.

Chapter 9

Finite Probability Distributions

Definition 9.1. A finite probability distribution is a function $f : A \rightarrow \mathbb{R}$ from a finite type A , such that f is nonnegative and the sum of f is 1.

Definition 9.2. The uniform distribution on a nonempty set A , $\text{Uniform}(A)$, assigns $\frac{1}{|A|}$ to all values in A and 0 to other values.

Definition 9.3. Given two finite probability distributions $f : A \rightarrow \mathbb{R}, g : B \rightarrow \mathbb{R}$, we have a probability distribution from $A \times B$ defines as $(f \times g)(x, y) = f(x)g(y)$.

Definition 9.4. Given a finite probability distribution $f : A \rightarrow \mathbb{R}$ and a function $g : A \rightarrow B$, we can apply g to the random variable represented by f . This gives the distribution $g\#f$.

We can directly transfer all theorems on $f\#g$ to finite PMFs.

Definition 9.5. Given two finite probability distributions $f : A \rightarrow \mathbb{R}, g : A \rightarrow \mathbb{R}$, we have a probability distribution defines as $f - g = s\#(f \times g)$ with $s(x, y) = x - y$.

Definition 9.6. Given two finite probability distributions $f : A \rightarrow \mathbb{R}, g : A \rightarrow \mathbb{R}$, we have a probability distribution defines as $f + g = a\#(f \times g)$ with $a(x, y) = x + y$.

Definition 9.7. Given a finite probability distribution $f : A \rightarrow \mathbb{R}$, we have a probability distribution defines as $-f = n\#f$ with $n(x) = -x$.

Proposition 9.8. These operations define a commutative monoid.

Lemma 9.9. We have $(f\#a) \times (g\#b) = h\#(a \times b)$, with $h(x, y) = (f(x), g(y))$.

Proof. By calculation. □

Lemma 9.10. We have $f\#(a \times b) = b \times a$ for $f(x, y) = (y, x)$.

Proof. Simple application of 8.5 □

Lemma 9.11. We have $(f\#a) + (g\#b) = h\#(a \times b)$, with $h(x, y) = f(x) + g(y)$.

Proof. By simplification after 9.9. □

Definition 9.12. Given a finite probability distribution $f : A \rightarrow \mathbb{R}$ and a list of finite probability distributions on B , indexed by elements of A , g , we can define $g(f)$ as the probability distribution obtained by sampling an element from f , and then sampling an element from the corresponding distribution in g .

Lemma 9.13. We have $f(g(a)) = h(a)$ with $h(x) = g(f(x))$.

Proof. By calculation. □

Lemma 9.14. We have $g\#f(a) = h(a)$ with $h(x) = g\#f(x)$.

Proof. By calculation. □

Definition 9.15. We say that a distribution a is ε -close to N entropy if for all sets $|A| \leq N$, $\sum_{x \in A} a(x) \leq \varepsilon$. Note that this is a bit different than the usual definition.

Proposition 9.16. If a is ε -close to $\lfloor n \rfloor$ entropy it's also ε -close to n entropy.

Proposition 9.17. If a is ε_1 -close to n entropy and $\varepsilon_1 \leq \varepsilon_2$ it's also ε_2 -close to n entropy.

Lemma 9.18. If e is an isomorphism and a is ε -close to n entropy, $e\#a$ is also ε -close to n entropy.

Proof. By definition, after using 8.5. □

Lemma 9.19. If a is ε -close to n entropy, then for any PMF b , $a+b$ is also ε -close to n entropy.

Proof.

$$\sum_{x \in H} (a+b)(x) = \sum_{x \in H} \sum_v b(v) a(x-v) = \sum_v b(v) \sum_{x \in H} a(x-v) = \sum_v b(v) \sum_{x \in H-v} a(x) \leq \sum_v b(v) \varepsilon = \varepsilon$$

□

Proposition 9.20. If, for all x such that $0 < f(x)$, we have that $g(x)$ is ε -close to n entropy, then $g(f)$ is ε -close to n entropy.

Proposition 9.21. For any probability distribution a , there are at most n values such that $a(x) > 1/n$.

Chapter 10

Lemmas about LP Norm

Theorem 10.1. For a function f with domain A

$$\|f\|_{\ell^1} \leq \sqrt{|A|} \|f\|_{\ell^2}$$

Proof. This is a particular case of the Cauchy-Schwartz inequality. □

Lemma 10.2. For a function f with domain A

$$\|f\|_{\ell^p} = |A|^{1/p} \|f\|_{L^p}$$

Proof. Trivial from the definition of $\|\cdot\|_{L^p}$. □

Lemma 10.3.

$$\|f\|_{\ell^p} \leq |A|^{1/p} \|f\|_{\ell^\infty}$$

Proof.

$$\left(\sum_x |f(x)|^p\right)^{1/p} \leq \left(\sum_x \|f\|_{\ell^\infty}^p\right)^{1/p} = (|A| \|f\|_{\ell^\infty}^p)^{1/p} = |A|^{1/p} \|f\|_{\ell^\infty}$$

.

Lemma 10.4. Note that in this lemma $\langle f, g \rangle$ is $\sum_x f(x)g(x)$.

$$|\langle f, g \rangle| \leq \|f\|_{\ell^1} \|g\|_{\ell^\infty}$$

Proof.

$$\left|\sum_x f(x)g(x)\right| \leq \sum_x |f(x)g(x)| \leq \sum_x |f(x)| \|g\|_{\ell^\infty} = \|f\|_{\ell^1} \|g\|_{\ell^\infty}$$

Lemma 10.5.

$$\|a\|_{\ell^2} \leq \sqrt{\|a\|_{\ell^1} \|a\|_{\ell^\infty}}$$

Proof. Trivial with [10.4](#) and $\|a\|_{\ell^2} = \sqrt{\langle a, a \rangle}$ □

Chapter 11

XOR Lemma

Most of the material in here was taken from [Rao07].

Theorem 11.1. *For a function f with domain A ,*

$$\|f\|_{\ell^1} \leq |A|^{3/2} \|\hat{f}\|_{\ell^\infty}$$

Proof. By 10.1 we have $\|f\|_{\ell^1} \leq \sqrt{|A|} \|f\|_{\ell^2}$. Then using 10.2 this is $|A| \|f\|_{L^2}$. By Parseval's theorem, this is $|A| \|\hat{f}\|_{\ell^2}$. By 10.3, we have $\|\hat{f}\|_{\ell^2} \leq \sqrt{|A|} \|\hat{f}\|_{\ell^\infty}$, which combines to the desired conclusion. \square

Lemma 11.2. *This is a very slight generalization of Lemma 4.3 in [Rao07]:*

Let G, H be finite abelian groups. Let X be a function $G \rightarrow \mathbb{R}$ such that for every nontrivial character χ , $\hat{X}(\chi) \leq \frac{\varepsilon}{|G|}$ and let U be the function with constant value $E_x[X(x)]$. Let $\sigma : G \rightarrow H$ be a function such that for every character ϕ , we have $\|\widehat{\phi \circ \sigma}\|_{\ell^1} \leq \tau$. Then $\|\sigma \# X - \sigma \# U\|_{\ell^1} \leq \tau \varepsilon \sqrt{|H|}$.

Proof. The proof is identical to the proof in [Rao07], using 11.1. \square

Lemma 11.3. *If a, b, n are reals, b, n are positive, and $\frac{a}{b} \leq n$, then $\frac{a}{b} \leq \frac{a+1}{b+1/n}$.*

Proof. By direct calculation (alternatively, this can be seen as an instance of the mediant inequality). \square

Lemma 11.4. *For a real x , we have $2 - |4x - 2| \leq |e^{x2\pi i} - 1|$.*

Proof. We have $|e^{x2\pi i} - 1| = |\cos(2\pi x) - 1 + i \sin(2\pi x)| = \sqrt{(\cos(2\pi x) - 1)^2 + \sin^2(2\pi x)} = \sqrt{2 - 2\cos(2\pi x)}$. WLOG, it's sufficient to consider the range $0 \leq x \leq \frac{1}{2}$. In this range, we have the inequality $\cos(2\pi x) \leq 1 - \frac{2}{\pi^2}(2\pi x)^2 = 1 - 8x^2$, from which the result quickly follows. \square

In the following, we consider $\sigma : \mathbb{Z}_N \rightarrow \mathbb{Z}_M$ defined as $\sigma(x) = x \bmod M$.

Lemma 11.5. *We have $\|\sigma \# U - U\|_{\ell^1} \leq \frac{n}{m}$.*

Proof. We can easily bound each difference by $\frac{1}{n}$ using $(\sigma \# U)(x) = \frac{\lceil \frac{N - (x \bmod M)}{M} \rceil}{N}$ and $U(x) = \frac{N}{N}$. \square

Theorem 11.6. *This is Lemma 4.4 in [Rao07] with explicit constants:*

For any character χ of \mathbb{Z}_M , $\|\widehat{\chi \circ \sigma}\|_{\ell^1} \leq 6 \ln(N) + 6$

Proof. Let $\rho(x) = e^{x2\pi i}$. We can find a value w such that $\chi(x) = \rho(wx/M)$. Then $\chi(\tau(x)) = \rho(wx/M)$. Now we have

$$\|\widehat{\chi \circ \sigma}\|_{\ell^1} = \frac{1}{N} \sum_{t \in \mathbb{Z}_N} \left| \sum_{x \in \mathbb{Z}_N} \rho(wx/M) \rho(-tx/N) \right| = \frac{1}{N} \sum_{t \in \mathbb{Z}_N} \left| \sum_{x \in \mathbb{Z}_N} \rho\left(\frac{wN-tM}{NM}\right)^x \right|$$

We now want to claim $\left| \sum_{x \in \mathbb{Z}_N} \rho\left(\frac{wN-tM}{NM}\right)^x \right| \leq \frac{|\rho\left(\frac{wN-tM}{NM}\right)^N - 1| + 1}{|\rho\left(\frac{wN-tM}{NM}\right) - 1| + 1/N}$. If $\rho\left(\frac{wN-tM}{NM}\right) = 1$, this is trivially correct. Otherwise, this is a geometric sum, and then we can use 11.3. We easily have $|\rho\left(\frac{wN-tM}{NM}\right)^N - 1| + 1 \leq 3$, and now we need to bound $\frac{1}{N} \sum_{t \in \mathbb{Z}_N} \frac{1}{|\rho\left(\frac{wN-tM}{NM}\right) - 1| + 1/N} = \frac{1}{N} \sum_{t \in \mathbb{Z}_N} \frac{1}{|\rho\left(\frac{wN-tM}{NM}\right) - 1| + 1/N}$. We can use 11.4 to bound this as $\frac{1}{N} \sum_{t \in \mathbb{Z}_N} \frac{1}{(2 - |4(\frac{wN-tM}{NM}) - 2|) + 1/N}$. By writing $wN/M = \lfloor wN/M \rfloor + \langle wN/M \rangle$, this is equal to $\sum_{t \in \mathbb{Z}_N} \frac{1}{2N - |4(\langle wN/M \rangle + t) - 2N| + 1}$. Now by splitting to cases and calculating we can see that $\frac{1}{2N - |4(\langle wN/M \rangle + t) - 2N| + 1} \leq \frac{1}{4t+1} + \frac{1}{4(n-1-t)+1}$. Applying bounds on the harmonic sum, we get the desired result. \square

Theorem 11.7. *Let X be a distribution \mathbb{Z}_N such that for every nontrivial character χ , $\hat{X}(\chi) \leq \frac{\varepsilon}{|G|}$. Then $\text{SD}(\sigma \# X, U) \leq \varepsilon \sqrt{M}(3 \ln(N) + 3) + \frac{M}{2N}$.*

Proof. Trivial with $\text{SD}(A, B) = \|A - B\|_{\ell^1}$, the triangle inequality with 11.5, 11.2 and 11.6. \square

[Rao07]: Rao, Anup. ‘‘An Exposition of Bourgain’s 2-Source Extractor.’’ Electron. Colloquium Comput. Complex. TR07 (2007): n. pag.

Chapter 12

Lemmas about the Inner Product Extractor

Proposition 12.1. For a character χ , $\chi(a) = \chi(b)$ iff $\chi(a - b) = 1$.

Proposition 12.2. The inner product is commutitive.

Lemma 12.3. If χ is a non-trivial character of a field \mathbb{F} , then there is an injective function from elements of \mathbb{F}^2 (generalize this to any dimension) to characters of it, defined by $f(x)(y) = \chi(x \cdot y)$.

Proof. It's easy to see this maps values to additive characters. For injectivity, we have some value x such that $\chi(x) \neq 1$. Now if $f((a_1, a_2)) = f((b_1, b_2))$, if they aren't equal, we can apply either $\frac{x}{a_1 - b_1}$ or $\frac{x}{a_2 - b_2}$, and then we get $\chi(x) = 1$ by 12.1, a contradiction. \square

Lemma 12.4. The function in the previous lemma is actually a bijection.

Proof. By 12.3 and the cardinality being equal. \square

Theorem 12.5.

Note: the inner product and DFT here aren't normalized.

$$\sum_x a(x) \sum_y b(y) \chi(x \cdot y) = \langle a, P(\hat{b}) \rangle$$

where P reorders \hat{b} based on 12.4

Proof. TODO \square

Theorem 12.6.

$$\left| \sum_x a(x) \sum_y b(y) \chi(x \cdot y) \right|^2 \leq |A|^2 \|a\|_{\ell^2}^2 \|b\|_{\ell^2}^2$$

Proof. We use 12.5 to rewrite the sum, and then use Cauchy-Schwartz. Then we can undo the reordering and use Parseval's theorem to get the desired result. \square

Theorem 12.7.

$$\left| \sum_x a(x) \sum_y b(y) \chi(x \cdot y) \right| \leq |A| \|a\|_{\ell^2} \|b\|_{\ell^2}$$

Proof. Simplying apply a square root to 12.6. \square

Theorem 12.8. For any bilinear form F and character χ ,

$$\left| \sum_x a(x) \sum_y b(y) \chi(F(x, y)) \right|^2 \leq \left| \sum_x a(x) \sum_y (b - b)(y) \chi(F(x, y)) \right|^2$$

Proof.

$$\left| \sum_x a(x) \sum_y b(y) \chi(F(x, y)) \right|^2 \leq \left(\sum_x a(x) \left| \sum_y b(y) \chi(F(x, y)) \right| \right)^2 \quad (12.1)$$

$$\leq \sum_x a(x) \left| \sum_y b(y) \chi(F(x, y)) \right|^2 \quad (12.2)$$

$$= \sum_x a(x) \left(\sum_y b(y) \chi(F(x, y)) \right) \left(\sum_y \overline{b(y) \chi(F(x, y))} \right) \quad (12.3)$$

$$= \sum_x a(x) \sum_y \sum_{y'} b(y) b(y') \chi(F(x, y)) \chi(-F(x, y')) \quad (12.4)$$

$$= \sum_x a(x) \sum_y \sum_{y'} b(y) b(y') \chi(F(x, y - y')) \quad (12.5)$$

$$= \sum_x a(x) \sum_y (b - b)(y) \chi(F(x, y)) \quad (12.6)$$

$$(12.7)$$

□

Theorem 12.9. For any bilinear form F and character χ ,

$$\left| \sum_x a(x) \sum_y b(y) \chi(F(x, y)) \right| \leq \sqrt{\left| \sum_x a(x) \sum_y (b - b)(y) \chi(F(x, y)) \right|^2}$$

Proof. Trivial from 12.8. □

Theorem 12.10. If a and b are ε -close to N entropy, then

$$\left| \sum_x a(x) \sum_y b(y) \chi(F(x, y)) \right| \leq \frac{|A|}{N} + 2\varepsilon$$

Proof. From the hypothesis and 9.21 we can look at $a'(x) = \begin{cases} a(x) & a(x) \leq \frac{1}{N} \\ 0 & \frac{1}{N} < a(x) \end{cases}$, and similarly for b' , and the difference would be at most 2ε . Then we can apply 12.7 to get the result. □

Chapter 13

Bourgain Extractor

Definition 13.1. Given a distribution A on \mathbb{F} , and a distribution B on \mathbb{F}^3 , we define a distribution $L(A, B)$ by sampling x from A , sampling (y, z, w) from B , and outputting $(x+y, z(x+y)+w)$.

Lemma 13.2. We have $L(f(A), g(B)) = L'(A \times B)$ with $L'(x, y) = L(f(x), g(y))$.

Proof. Trivial with ?? and 9.14. □

Theorem 13.3. Given an integer N and a real number β such that $p^\beta \leq N \leq p^{2-\beta}$, and two nonempty sets $A' \subseteq \mathbb{F}, B' \subseteq \mathbb{F}^3$, such that $|B'| \leq N$ and the last two values in every element of B' are unique, then $L(\text{Uniform}(A'), \text{Uniform}(B'))$ is $\frac{C}{|A'| |B'|} N^{3/2-\varepsilon(\beta)}$ -close to N entropy.

Proof. TODO □

Theorem 13.4. TODO

Proof. TODO □

Theorem 13.5. TODO

Proof. TODO □

Definition 13.6. $M(x, y) = (x + y, 2(x + y), -((x + y)^2 + x^2 + y^2))$

Definition 13.7. $D(x, y) = (x, x^2 - y)$.

Lemma 13.8. $f\#(b \times b \times b) = D\#L(b, M\#(b \times b))$, with $f(x, y, z) = (x + y + z, x^2 + y^2 + z^2)$.

Proof. By direct calculation with 9.9, 8.6, 9.10. □

Lemma 13.9. If the maximum value of a is ε , the maximum value of $M\#(a \times a)$ is at most $2\varepsilon^2$.

Proof. It suffices to show that every value can be obtained at most twice as an output of M . Because the first value determines the second one, we can drop it, and then if want to get (x_1, x_2) we need $y_1 + y_2 = x_1, y_1 y_2 = x_1^2 + x_2/2$ (by calculation). A calculation can further show that $(x_1, x_2) \rightarrow -x_1$ is a bijection from this to the set of roots of $y^2 + x_1 y + (x_1^2 + x_2/2)$, which is easily of size at most 2. □

Definition 13.10. $\beta = \frac{35686629198734976}{35686629198734977}$.

Definition 13.11. $\alpha = \varepsilon(\beta)$

Lemma 13.12. $\alpha = \frac{11}{2}(1 - \beta)$.

Proof. By calculation. \square

Lemma 13.13. For any source a with maximum value at most $p^{-1/2+2/11\alpha}$, $D\#L(a, M\#(a \times a))$ is $8Cp^{-2/11\alpha}$ -close to $p^{1+2/11\alpha}$ entropy.

Proof. First, by 9.18, we can get rid of the D . Now we want to apply 13.5. We already have a bound for the maximum value of a , and using 13.9 we get a bound for the maximum value of $M\#(a \times a)$. The last two values of a triple in the support $M\#(a \times a)$ is an injective function by 8.10, as the first value is half of the second value for triples in the domain of M . \square

Definition 13.14. $C_b = \sqrt[64]{16C + 1}$.

Theorem 13.15. For any two sources a, b with maximum value at most $p^{-1/2+2/11\alpha}$, and any non-trivial character χ ,

$$\left| \sum_x a(x) \sum_y b(y) \chi(xy + x^2y^2) \right| \leq C_b p^{-1/352\alpha}$$

Proof. First define $a' = f\#a, b' = f\#b$ for $f(x) = (x, x^2)$, then this is $|\sum_x a'(x) \sum_y b'(y) \chi(x \cdot y)|$. Applying 12.9 3 times, then swapping x, y and doing it three more times, we can bound this by $|\sum_x (b' + b' + b' + (b' - b' - b' - b' - b'))(x) \sum_y (a' + a' + a' + (a' - a' - a' - a' - a'))(y) \chi(x \cdot y)|^{1/64}$. Now we want to use 12.10. By 9.19, it suffices to show that $b' + b' + b'$ and $a' + a' + a'$ are close to high entropy. First, we can rewrite this by unfolding a' and b' , using 9.11 and then 13.8. Finally, what we want is 13.13. \square

Theorem 13.16. For any positive integer m and two sources a, b with maximum value at most $p^{-1/2+2/11\alpha}$, the statistical distance of $f\#(a \times b)$ with $f(x, y) = (xy + x^2y^2 \bmod p) \bmod m$ to the uniform distribution is at most $\varepsilon = C_b p^{-1/352\alpha} \sqrt{m}(3 \ln(p) + 3) + \frac{m}{2p}$.

Proof. This is a simple application of 11.7 with 13.15 \square

Theorem 13.17. For any positive integer m , the function $f(x, y) = (xy + x^2y^2 \bmod p) \bmod m$ is a two source extractor, with $k = (1/2 - 2/11\alpha) \log(p)$, $\varepsilon = C_b p^{-1/352\alpha} \sqrt{m}(3 \ln(p) + 3) + \frac{m}{2p}$.

Proof. This is a simple restatement of 13.16 \square